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# IsaPlanner: A Proof Planner for Isabelle

Lucas Dixon

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# Motivation

- Improve automation in Isabelle by using tools from proof planning
- A generic approach?
- Take advantage of the proof tools available in Isabelle
- Expressing the reasoning process and the resulting proof plans in a *readable* form (Surveyable proofs)
- Proof planning to aid the interactive user

# IsaPlanner Approach to Proof Planning

- Encode Techniques: capture common patterns of reasoning
- Applying a technique generates high level descriptions of the proofs called *proof plans* (a representation of proof scripts)
- Our approach derives Isar style proof scripts.
- Rippling: a difference removal technique, principally used to guide rewriting of the step case during inductive proofs

## IsaPlanner

- A framework for encoding and applying reasoning techniques that derives readable proof scripts.
- Mechanisms for sharing non-logical information between proof techniques.
- Applying proof techniques produces Isar proof scripts
  - Allow human inspection of the proof attempt
  - Support debugging of proof
  - Support the exploration of proof
- Provide well-behaved techniques

# Hierarchy of Behaviour

1. Bug
2. Non termination
3. Uninformative result (big goals, exceptions, "no")
4. Informative result (meaningful progress)
5. Proof/Refutation

## Example Generated Proof Script

```
theorem sum_of_odds:  $(\sum_{i < n} 2 * i + 1) = n^2$ 
  proof (induct n)
    show  $(\sum_{i < 0} 2 * i + 1) = 0^2$  by simp
  next
    fix n
    assume IH:  $(\sum_{i < n} 2 * i + 1) = n^2$ 
    have  $(\sum_{i < n} 2 * i + 1) + ((2 * n) + 1) = (n + 1)^2$ 
      gap (conjecture_lemma  $n^2 + ((2 * n) + 1) = (n + 1)^2$ )
    thus  $(\sum_{i < (n + 1)} 2 * i + 1) = (n + 1)^2$  by rippling
  qed
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```

## Proof Planning?

- Declarative description (datatype) of Isar proof scripts
- Compared with tactics:
  - More information - history of the proof
  - Surveyable proofs
- Mix human/computer manipulation of 'readable' formal proof



## Encoded Techniques

- Induction, Case splitting, Equational Reasoning
- Higher order rippling (provides rewriting)
- Lemma speculation and generalisation critics
- This combination forms a powerful inductive theorem prover

## Some Observations

- Supports *exploration* of Isar proofs
- Can prove some apparently complicated problems, for example:  $(x^y)^z = x^{(y \cdot z)}$  in ordinal arithmetic.
- Conjecturing lemmas is a directed form of theory formation
- Unsuccessful branches/failed proof attempts can still result in proving useful lemmas
- Caching conjectures (proved and unprovable) prunes the search space significantly

## Some Issues

- Where to draw the line between work done by a tactic and work done by proof planning?
- What information to store in source files for proof re-execution?
- Interface issues: how to use support interactive proof planning? (Proof by contextual menu!)