

# Formalising Lagrange's Theorem in Coq

---

Dan Synek

Radboud University Nijmegen

*WHY?*

## Informal Proof

---

Let  $G$  be a finite group. The order of a subgroup  $H$  of  $G$  divides the order of  $G$ .

Informal Proof: Form the cosets  $Hg$  of  $G$  where  $Hg := \{hg \mid h \in H\}$

1. The cosets form a partition, e.g.  $Hg_1 = Hg_2$  or  $Hg_1 \cap Hg_2 = \emptyset$ :  
Assume  $Hg_1 \cap Hg_2 \neq \emptyset$  so  $h_1g_1 = h_2g_2$  and thus  $g_1 = h_1^{-1}h_2g_2$ ,  
then  $g_1 \in Hg_2$  and thus  $Hg_1 \subseteq Hg_2$ . By symmetry  $Hg_2 = Hg_1$ .
2. Each coset has exactly  $|H|$  elements: If  $h_1g = h_2g$  then  $h_1 = h_2$ .

## Informal Version of Formal Proof

---

1.  $G \simeq G/H \times H$
2.  $|A \times B| = |A||B|$

where the equality on  $G/H$  is defined by  $\hat{x} = \hat{y}$  iff  $xy^{-1} \in H$ . We do not always have:

$$\forall a : A \exists b : B, (Rab) \rightarrow \exists f : (A \Rightarrow B), \forall a : A, (Ra(fa))$$

but we can prove it for finite  $A$ .

The isomorphism is given by:

$$f(g) = \langle \hat{g}, g \cdot (c(g))^{-1} \rangle$$

$$f^{-1} \langle \hat{g}, h \rangle = h \cdot (c(g))$$

(2) is proved by induction over the size of  $B$ .

## What I had and what needed to be done

---

Corn has:

- Setoids, including functions, union and products between setoids.
- Groups, including simplification of algebraic expressions.

I had to do:

- Define finiteness.
- Define isomorphisms and prove lemmas about it.
  - $\simeq$  forms equivalence,  $A \simeq (S \mid P) \cup (A \mid \neg P)$  and  $A \times (B \cup C) \simeq A \times B \cup A \times C$  about 15 more
- Relate decidability to isomorphisms, finiteness and subsets.
- Define extensional choice and prove it for finite sets.

## So what is the problem with Setoids

---

A Setoid is a type and an equivalence relation.

- No obvious choice how to represent set operations. Should one work with predicates representing sets, with union as predicate operation, or work with disjoint union on the carriers? Is a subgroup and group based on different notion of set?
- Intensionality of type theory together with setoids often forces you to use isomorphisms, e.g. subset and quotient don't commute.
- A common construction is to make a setoid based on a sigma type with the carrier and a predicate, with the equality "inherited" from the original setoid. In practice using these setoids leads to extensive packing and unpacking (as observed by Luis Cruz-Filipe).
- $A \mid P \mid Q \simeq A \mid Q \mid P$  is not a meaningful statement.

## Some observations

---

- Around 6000 lines of proof.
- Most lines are spent on proving lemmas about isomorphisms.
- We could hide the size of the proof behind fancy tactics. But is that what we want?
- Things like better syntax, setoid rewriting and a big library is not the solution.

## Future Work

---

- Formalising can be a nice way to illustrate the power of abstraction. Motivation for proof assistants as a tool in teaching.
- Is PER/subsetoids a better framework to represent sets?
- Or more radically: Work in ZF coded in type theory a la Miguel
- Or even more radically: Find an application that justifies the choice of framework!